

# Rank Aggregation by Dissatisfaction Minimisation in the Unavailable Candidate Model

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## Introduction

Several approaches have recently been proposed to address the problem of candidates' unavailability in social choice, among which the *unavailable candidate model* proposed in [1] where the optimal rankings are computed by minimisation of the expected number of disagreements over all the possible subsets of available candidates. Lu and Boutilier [1] provide a clear justification for producing a ranking instead of a single winner: the output ranking serves as a very compact decision policy to select the best available candidate as "winner".

- $C$ : set of all the candidates
- $top_r(S)$ : best ranked candidate in subset  $S$ , for ranking (or voter)  $r$
- $p \in [0,1]$ : probability for any candidate to be unavailable, leading to a Bernoulli distribution over the subsets of  $C$
- The optimal rankings minimise the sum of dissatisfaction among all the voters, where the dissatisfaction between  $r$  and  $r'$  is

$$D_p(r, r') = \sum_{S \subseteq C} p^{m-|S|} (1-p)^{|S|} \mathbf{1}[top_r(S) \neq top_{r'}(S)]$$

## Generalized unavailable candidate model

- Binary disagreement in [1] -> a voter is satisfied if her preferred available candidate is elected and fully unsatisfied otherwise.
- In our work, the voter's satisfaction varies more smoothly and depends on the rank she gives to the winner.
- Two opposed ways to measure the satisfaction of the voters. In the following,
- $\rho$  measures the disutility associated with a rank
- $P$ : the distribution of probability on the sets of available candidates

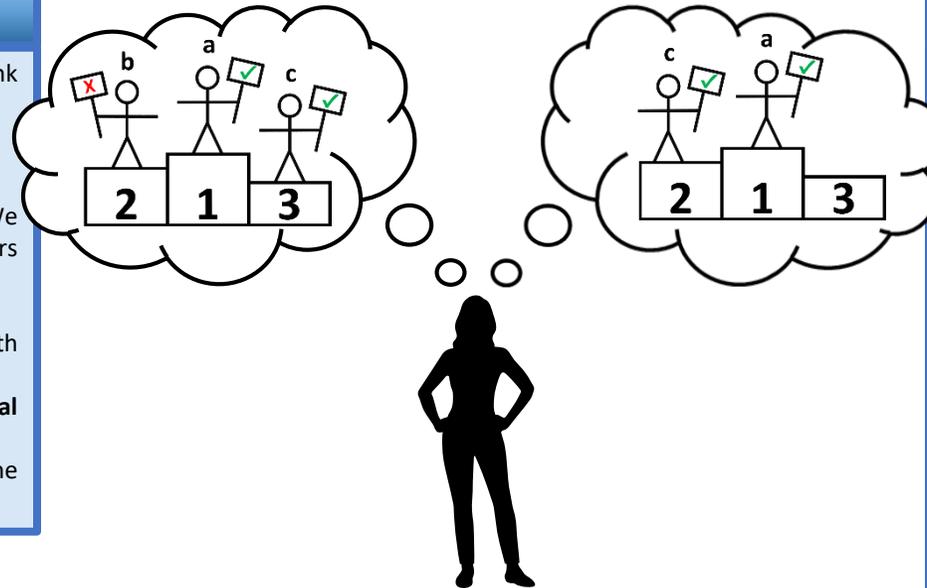
## Ex-ante approach

The disutility felt by the voter  $v$  is computed considering the rank of the candidates among all the candidates of  $C$ :

$$\hat{\Delta}_{\rho, P}(v, r) := E_{S \sim P} [\rho(v(top_r(S))) - \rho(v(top_v(S)))]$$

**Ex. 1:** Let  $C=\{a,b,c\}$ ,  $\rho=(0,1,2)$ ,  $P$  the uniform distribution. We consider 4 voters with the preference order  $(a,b,c)$  and 7 voters with the order  $(c,a,b)$ . The only optimal ranking is **(a,c,b)**.

- The ex-ante dissatisfaction rule amounts to a **scoring rule** with score  $-\rho$ 
  - finding an optimal ranking can be done in **polynomial time**
  - the optimal rankings do not depend on the distribution  $P$



## Ex-post approach

The disutility felt by the voter  $v$  is computed considering the rank of the candidates *within* the set of available candidates:

$$\Delta_{\rho, P}(v, r) := E_{S \sim P} [\rho(v_S(top_r(S))) - \rho(v_S(top_v(S)))]$$

**Ex. 2:** With the same setting as for Ex. 1, the only optimal ranking is now **(c,a,b)**.

- Under reasonable assumptions on  $P$  and  $\rho$ , any optimal ranking is also a **Kemeny consensus**
  - the ex-post dissatisfaction rule **NP-hard**.
- We exhibit a **polynomial-time approximation scheme**
  - the ex-post rule can be handled in practise.
- Outline of the algorithm:
  - ❖ While there is a dominant candidate (satisfying a property ensuring it is first in any optimal ranking) append it to the output ranking and remove it from  $C$
  - ❖ Choose and order  $K$  candidates in  $C$  such that their contribution to  $\Delta_{\rho, P}(v, r)$  is minimal
  - ❖ Randomly order the remaining candidates.

## Voting rules covered by our model

Scoring rules, positional Spearman distance.

Kemeny rule, rules that minimise the expected number of disagreements under different probability distributions ([1], [2], [3]).

[1] Tyler Lu and Craig E Boutilier. 2010. The unavailable candidate model: a decision-theoretic view of social choice. In *Proceedings of the 11th ACM conference on Electronic commerce*. 263–274.

[2] Katherine A Baldiga and Jerry R Green. 2013. Assent-maximizing social choice. *Social Choice and Welfare* 40, 2 (2013), 439–460.

[3] Hugo Gilbert, Tom Portoleau, and Olivier Spanjaard. 2020. Beyond Pairwise Comparisons in Social Choice: A Setwise Kemeny Aggregation Problem. In *Thirty-Fourth AAAI Conference on Artificial Intelligence*. 1982–1989.